

Last time

Numerical sensitivity of basic operations

Basic representations (why we won't deal with it directly)

Today - start with numerical linear algebra

Solving $Ax=b$

- vector & matrix norms
- condition number $K(A)$
- direct methods
 - Gaussian elimination
 - LU decomposition
 - Cholesky decomposition

Norms (vector)

$$\|\cdot\| : \mathbb{C}^n \rightarrow \mathbb{R} \quad (\text{why } \mathbb{C}?)$$

→ difference between \mathbb{C} & \mathbb{R}

satisfies

positive : $\|x\| \geq 0 \quad \forall x \in \mathbb{C}^n$

definiteness : $\|x\| = 0 \Leftrightarrow x = 0$

homogeneity : $\|\alpha x\| = |\alpha| \|x\| \quad \forall \alpha \in \mathbb{C} \quad \& \quad x \in \mathbb{C}^n$

triangle inequality : $\|x+y\| \leq \|x\| + \|y\|$

Example

$$x = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \quad x_i \in \mathbb{C} \Rightarrow x \in \mathbb{C}^n$$

p-norm

$$\|x\|_p = \left(\sum_{i=1}^n |x_i|^p \right)^{1/p} \quad \text{popular } p = 1, 2, \infty$$

(in probability)
3, 4, 2+ε

$p = \infty$ (max or sup norm)

$$\|x\|_\infty = \max \{ |x_1|, |x_2|, \dots, |x_n| \}$$

Idea $\lim_{p \rightarrow \infty} \|x\|_p \rightarrow \max \{ |x_i| \}$

$$M = \max \{ |x_i| \} \text{ so } |x_i| \leq M \quad \forall i$$

$$\left(\sum_{i=1}^n |x_i|^p \right)^{1/p} \leq \left(\sum_{i=1}^n M^p \right)^{1/p} = (n M^p)^{1/p} = n^{1/p} M$$

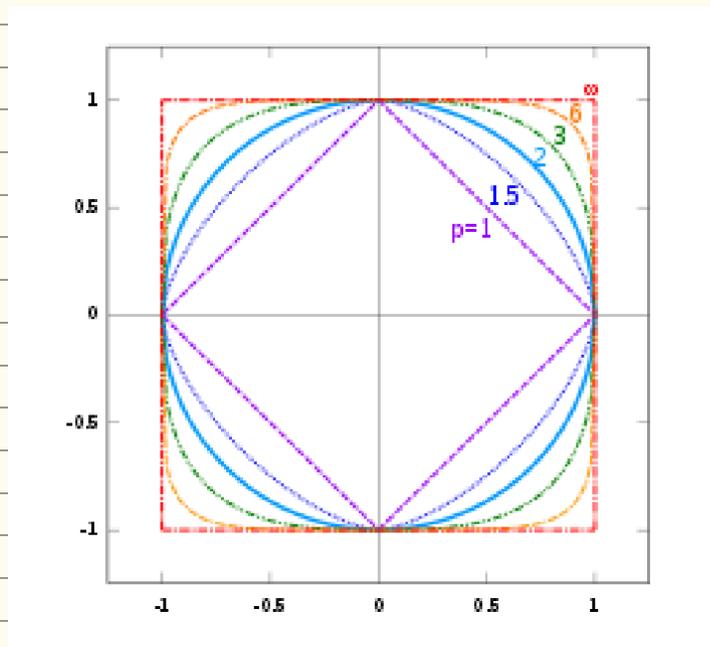
since there is an i st $|x_i| = M$

$$\left(\sum_{i=1}^n |x_i|^p \right)^{1/p} \geq (M^p)^{1/p}$$

$$M \leq \|x\|_p \leq n^{1/p} M \quad p \rightarrow \infty$$

$$n^{1/p} \rightarrow 1 \quad \text{if } n \geq 1$$

Geometric visualization of norms



Note: norms for $p < 1$ are not convex
& $p=0$ is usually defined as
of non zero entries

Matrix norms

$$\|\cdot\|: \mathbb{C}^{n \times n} \rightarrow \mathbb{R}$$

• positive definiteness $\|A\| \geq 0 \quad \forall A \in \mathbb{C}^{n \times n}$
 $\|A\| = 0 \Leftrightarrow A = 0$

• homogeneity $\|\alpha A\| = |\alpha| \cdot \|A\| \quad \forall \alpha \in \mathbb{C}$
 $A \in \mathbb{C}^{n \times n}$

• triangle inequality $\|A+B\| \leq \|A\| + \|B\|$
 $\forall A, B \in \mathbb{C}^{n \times n}$

• submultiplicativity $\|A \cdot B\| \leq \|A\| \cdot \|B\| \quad \forall A, B \in \mathbb{C}^{n \times n}$

Claim :

For every vector norm $\|\cdot\|_*$ on \mathbb{C}^n

$$\|A\|_* = \max_{\|x\|=1} \|Ax\|_* = \max_{\|x\| \neq 0} \frac{\|Ax\|_*}{\|x\|_*}$$

determines a matrix norm on $\mathbb{C}^{n \times n}$

Proof : Positive : if $A \neq 0 \exists x' \in \mathbb{C}^n$ st.

Definiteness $\|x'\|=1 \quad \& \quad Ax' \neq 0$

$$\text{so } \|A\|_* \geq \|Ax'\|_* \geq \|Ax'\|_* > 0$$

$$\text{if } A=0 \quad Ax=0 \Rightarrow \|A\|_* = 0$$

homogeneity : $\forall \alpha \quad \|\alpha Ax\|_* = |\alpha| \|Ax\|_*$

$$\|\alpha A\|_* = \max_{\|x\|=1} \|\alpha Ax\|_* =$$

$$= \max_{\|x\|=1} |\alpha| \|Ax\|_*$$

$$= |\alpha| \max_{\|x\|=1} \|Ax\|_* = |\alpha| \|A\|_*$$

Triangle inequality

$$\|A+B\|_* = \max_{\|x\|=1} \|(A+B)x\|_*$$

$$= \max_{\|x\|=1} \|Ax+Bx\|_*$$

$$\leq \max_{\|x\|=1} \|Ax\|_* + \max_{\|x\|=1} \|Bx\|_*$$

$$\leq \max_{\|x\|=1} \|Ax\|_* + \max_{\|y\|=1} \|By\|_*$$

$$= \|A\|_* + \|B\|_*$$

Submultiplicativity

$$\|A \cdot x\| \leq \|A\| \|x\|$$

$$\|A \cdot B\|_* = \max_{\|x\|=1} \frac{\|ABx\|_*}{\|x\|_*}$$

$$= \max_{\|x\|_* \neq 0, \|Bx\|_* \neq 0} \left(\frac{\|ABx\|_*}{\|Bx\|_*} \cdot \frac{\|Bx\|_*}{\|x\|_*} \right)$$

$$\leq \max_{\|Bx\|_* \neq 0} \frac{\|ABx\|_*}{\|Bx\|_*} \cdot \max_{\|x\|_* \neq 0} \frac{\|Bx\|_*}{\|x\|_*}$$

$$\text{ce } \|Bx\|_* = 0 \rightsquigarrow \leq \max_{y \neq 0} \frac{\|Ay\|_*}{\|y\|_*} \cdot \max_{\|x\|_* \neq 0} \frac{\|Bx\|_*}{\|x\|_*}$$

$$\|AB\|_* = 0$$

$$= \|A\|_* \cdot \|B\|_*$$

Examples

1-norm $\|A\|_1 = \max_j \left(\sum_{i=1}^n |a_{ij}| \right)$ (Proof?)

2-norm (spectral norm)

$$\|A\|_2 = \sqrt{\max_j \lambda_j(A^T A)}$$

Frobenius norm

$$\|A\|_F = \left(\sum_{i,j} |a_{ij}|^2 \right)^{1/2}$$

Supremum norm

$$\|A\|_\infty = \max_i \left(\sum_j |a_{ij}| \right)$$

Nuclear norm (Schatten norm)

$$\|A\|_* = \sum_j \sqrt{\lambda_j(A^T A)}$$

Why different norms

- Computation different complexity
- all "equivalent" (to some extent)

$$\frac{1}{\sqrt{n}} \|A\|_F \leq \|A\|_2 \leq \|A\|_F$$

$$\frac{1}{\sqrt{n}} \|A\|_1 \leq \|A\|_2 \leq \sqrt{n} \|A\|_1$$

$$\frac{1}{\sqrt{n}} \|A\|_\infty \leq \|A\|_2 \leq \sqrt{n} \|A\|_\infty$$

Can be good for fixed n (dimension)

Sensitivity of $Ax=b$

Question: Given $Ax=b$, how does the solution x change if we perturb A & / or b

$$(A + \Delta A)(x + \Delta x) = b + \Delta b$$

Goal: Estimate $\frac{\|\Delta x\|}{\|x\|} \rightsquigarrow$ vector norm

Sensitivity / condition number

$$\kappa(A) = \|A\| \cdot \|A^{-1}\|$$

Theorem

Let A be an invertible matrix (A^{-1} exists)

① if $\Delta A = 0$ (only care about Δb)

$$\frac{\|\Delta x\|}{\|x\|} \leq K(A) \frac{\|\Delta b\|}{\|b\|}$$

② if $\Delta A \neq 0$, let I denote the identity matrix $\begin{cases} \|A^{-1}\| \|\Delta A\| < 1 \end{cases}$

$$\frac{\|\Delta x\|}{\|x\|} \leq \frac{K(A)}{1 - K(A) \frac{\|\Delta A\|}{\|A\|}} \left(\frac{\|\Delta b\|}{\|b\|} + \frac{\|\Delta A\|}{\|A\|} \right)$$

Interpretation of ① - know system measurement error

Proof of ①

$$Ax = b \quad A\tilde{x} = \tilde{b} = b + \Delta b$$

\downarrow
 $x + \Delta x$

Upper bound

$$A(\tilde{x} - x) = \Delta b$$

$$A(x + \Delta x - x) = \Delta b$$

$$A\Delta x = \Delta b$$

$$\Delta x = A^{-1}\Delta b \Rightarrow \|\Delta x\| \leq \|A^{-1}\| \|\Delta b\|$$

Lower bound

$$Ax = b \Rightarrow \|b\| = \|Ax\| \leq \|A\| \cdot \|x\|$$

$$\|x\| \geq \frac{\|b\|}{\|A\|} \Leftrightarrow \frac{1}{\|x\|} \leq \frac{\|A\|}{\|b\|} \quad *$$

Combined

$$\frac{\|\Delta x\|}{\|x\|} \leq \frac{\|\Delta^{-1}\| \cdot \|\Delta b\|}{\|b\| / \|A\|}$$

upper bound
lower bound

$$= \|A\| \cdot \|\Delta^{-1}\| \cdot \frac{\|\Delta b\|}{\|b\|}$$

$$= \kappa(A) \frac{\|\Delta b\|}{\|b\|}$$

Proof of ②

$$(A + \Delta A)(x + \Delta x) = b + \Delta b$$

$$\underbrace{Ax} + \Delta Ax + A \Delta x + \Delta A \Delta x = \underbrace{b + \Delta b}$$

we can cancel them since $Ax = b$

$$(\Delta A)x + A \Delta x + (\Delta A)(\Delta x) = \Delta b$$

$$(A + \Delta A) \Delta x = \Delta b - (\Delta A)x$$

→ Assume $A + \Delta A$ is invertible

$$\Delta x = (A + \Delta A)^{-1} \cdot (\Delta b - (\Delta A)x)$$

→ take the norm

$$\begin{aligned} \|\Delta x\| &\leq \|(A + \Delta A)^{-1}\| \cdot \|\Delta b - (\Delta A)x\| \\ &\leq \|(A + \Delta A)^{-1}\| \cdot (\|\Delta b\| + \|(\Delta A)x\|) \quad \left. \begin{array}{l} \text{triangle} \\ \text{ineq} \end{array} \right\} \\ &\leq \|(A + \Delta A)^{-1}\| \cdot (\|\Delta b\| + \|(\Delta A)\| \cdot \|x\|) \end{aligned}$$

Using $\textcircled{*}$ $\frac{1}{\|x\|} \leq \frac{\|A\|}{\|b\|}$

$$\Rightarrow \frac{\|\Delta x\|}{\|x\|} \leq \|(A + \Delta A)^{-1}\| \cdot \left(\|A\| \cdot \frac{\|\Delta b\|}{\|b\|} + \|\Delta A\| \right) \textcircled{*}$$

Express in terms of $A \in K(A)$

$$A + \Delta A = A (I + A^{-1}(\Delta A)) \quad \text{recall } (xy)^{-1} = y^{-1}x^{-1}$$

$$\begin{aligned} \|(A + \Delta A)^{-1}\| &\leq \|(I + A^{-1}(\Delta A))^{-1} \cdot A^{-1}\| \\ &\leq \|(I + A^{-1}(\Delta A))^{-1}\| \cdot \|A^{-1}\| \quad (*) \end{aligned}$$

Use series expansion

$$\|(I + A)^{-1}\| \leq \frac{1}{1 - \|A\|} \quad \left(\begin{array}{l} \text{just like for} \\ \text{scalars} \end{array} \right)$$

$$\begin{aligned} \|(I + A^{-1}(\Delta A))^{-1}\| &\leq \frac{1}{1 - \|A^{-1}(\Delta A)\|} \\ &\leq \frac{1}{1 - \|A^{-1}\| \cdot \|\Delta A\|} \end{aligned}$$

Putting it together

$$\begin{aligned} \|(A + \Delta A)^{-1}\| &\leq \|(I + A^{-1}(\Delta A))^{-1}\| \|A^{-1}\| \\ &\leq \frac{\|A^{-1}\|}{1 - \|A^{-1}\| \cdot \|\Delta A\|} \end{aligned}$$

Insert into (*)

$$\frac{\|\Delta x\|}{\|x\|} \leq \frac{\|A^{-1}\|}{1 - \|A^{-1}\| \|\Delta A\|} \left(\|A\| \cdot \frac{\|\Delta b\|}{\|b\|} + \|\Delta A\| \right)$$

$$= \frac{K(A)}{1 - \|A^{-1}\| \|\Delta A\|} \left(\frac{\|\Delta b\|}{\|b\|} + \frac{\|\Delta A\|}{\|A\|} \right)$$

$$= \frac{K(A)}{1 - K(A) \cdot \frac{\|\Delta A\|}{\|A\|}} \cdot \left(\frac{\|\Delta b\|}{\|b\|} + \frac{\|\Delta A\|}{\|A\|} \right)$$

$$\|A^{-1}\| \cdot \|\Delta A\| = \|A^{-1}\| \cdot \|A\| \cdot \frac{\|\Delta A\|}{\|A\|}$$

Example 1

Line example from week 1

$$A_1 = \begin{pmatrix} 1 & 1 \\ -1 & 1 \end{pmatrix}$$

stable

$$A_2 = \begin{pmatrix} 1 & 0.99 \\ 0.99 & 0.98 \end{pmatrix}$$

unstable

$$\kappa_1(A_1) = \kappa_F(A_1) = \kappa_\infty(A_1) = 2$$

$$\kappa_2(A_1) = 1$$

$$\kappa_1(A_2) = \kappa_\infty(A_2) = 3.96 \cdot 10^4$$

$$\kappa_2(A_2) = 3.94 \cdot 10^4$$

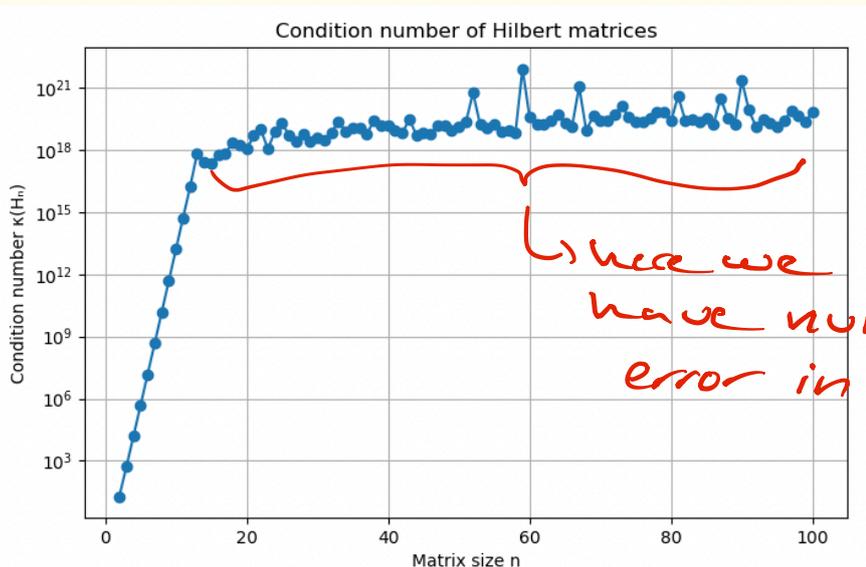
$$\kappa_F(A_2) = 3.92 \cdot 10^4$$

Example 2

Hilbert matrix $H_n = \left(\frac{1}{i+j-1} \right)_{ij} \in \mathbb{R}^{n \times n}$
($i, j = 1 \dots n$)

$$2 \times 2 \quad H_2 = \begin{pmatrix} 1 & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{3} \end{pmatrix}$$

(κ_2)



here we have numerical error in calculation

Direct Solution Methods

- Gaussian elimination
- LU factorization
- Pivoting (growth of pivots)
- Cholesky decomposition

Gaussian Elimination

Solving square linear systems

n equations, n unknowns

$$\begin{aligned} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n &= b_1 \\ a_{21}x_1 + \dots + a_{2n}x_n &= b_2 \\ \vdots & \\ a_{n1}x_1 + \dots + a_{nn}x_n &= b_n \end{aligned}$$

$$\begin{bmatrix} a_{11} & \dots & a_{1n} \\ \vdots & & \vdots \\ a_{n1} & \dots & a_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$$

A

x

If A is invertible $\left\{ \begin{bmatrix} a_{i1} \\ \vdots \\ a_{in} \end{bmatrix} \right\}_i$ forms a basis of \mathbb{R}^n

Look for coefficients such that

$$\begin{bmatrix} a_{11} \\ \vdots \\ a_{n1} \end{bmatrix} x_1 + \begin{bmatrix} a_{12} \\ \vdots \\ a_{n2} \end{bmatrix} x_2 + \dots + \begin{bmatrix} a_{1n} \\ \vdots \\ a_{nn} \end{bmatrix} x_n = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$$

given vector

coordinates of b in terms of $\left\{ \begin{bmatrix} a_{1i} \\ \vdots \\ a_{ni} \end{bmatrix} \right\}_i$

Cost of solving $Ax=b$

- computational complexity: # of operations
 $\{+, -, \cdot, : \}$

- error analysis (in terms of $K(A)$)

* For which matrices is the problem "easy" (for both complexity & error analysis)

Gaussian Elimination Review

$$w = 2, 2, \dots$$

Case 1: compute $A^{-1} \sim O(n^2)$ (or $O(n^3)$ practically)

We do not need A^{-1}

\rightarrow solve it "half way"

$$A = \begin{bmatrix} -3 & 2 & -1 \\ 6 & -6 & 7 \\ 3 & -4 & 4 \end{bmatrix} \quad b = \begin{bmatrix} -1 \\ -7 \\ -6 \end{bmatrix}$$

if we do row operations on A & b

$\Rightarrow x$ does not change

Idea: zero out columns below

$$A = \begin{bmatrix} -3 & 2 & -1 \\ 6 & -6 & 7 \\ 3 & -4 & 4 \end{bmatrix} \quad b = \begin{bmatrix} -1 \\ -7 \\ -6 \end{bmatrix}$$

⇓ write it this way

$$\left[\begin{array}{ccc|c} -3 & 2 & -1 & -1 \\ 6 & -6 & 7 & -7 \\ 3 & -4 & 4 & -6 \end{array} \right]$$

$$\begin{array}{l} x \\ y \end{array} \left[\begin{array}{ccc|c} -3 & 2 & -1 & -1 \\ 6 & -6 & 7 & -7 \\ 3 & -4 & 4 & -6 \end{array} \right] \rightarrow \begin{array}{l} x \\ z \end{array} \left[\begin{array}{ccc|c} -3 & 2 & -1 & -1 \\ 0 & -2 & 5 & -9 \\ 3 & -4 & 4 & -6 \end{array} \right]$$

$$y = y + 2x$$

$$z \leftarrow z + x$$

$$\left[\begin{array}{ccc|c} -3 & 2 & -1 & -1 \\ 0 & -2 & 5 & -9 \\ 0 & -2 & 3 & -7 \end{array} \right] \begin{array}{l} y \\ z \end{array} \left[\begin{array}{ccc|c} -3 & 2 & -1 & -1 \\ 0 & -2 & 5 & -9 \\ 0 & -2 & 3 & -7 \end{array} \right]$$

$$z \leftarrow z - y$$

$$\left[\begin{array}{ccc|c} -3 & 2 & -1 & -1 \\ 0 & -2 & 5 & -9 \\ 0 & 0 & -2 & 2 \end{array} \right] \rightsquigarrow \text{upper triangular}$$

$$-2x_3 = 2 \Rightarrow x_3 = -1$$

$$-2x_2 + 5x_3 = -9 \Rightarrow -2x_2 - 5 = -9 \Rightarrow x_2 = 2$$

$$-3x_1 + 2x_2 - x_3 = 2 \Rightarrow -3x_1 + 4 + 1 = 2 \Rightarrow x_1 = 1$$

Back substitution