

Fixed point approach

$$f(x) = 0 \Leftrightarrow g(x) = x$$

Alg: Choose x_0

$x_{k+1} = g(x_k)$ until convergence

Thm: If g is continuously differentiable on $I = [a, b]$

∴ $g(I) \subseteq I$ (closed under g) with

$$\sup_{x \in I} |g'(x)| = m < 1$$

① $g(x) = x$ has a unique solution ξ in I

② $\forall x_0 \in I$ the sequence $x_n = g(x_{n-1})$ converges to ξ w/

$$|x_{n+1} - \xi| \leq \min \left\{ m^{n+1} |x_0 - \xi|, \frac{m}{m-1} |x_{n+1} - x_n|, \frac{m^{n+1}}{1-m} |x_1 - x_0| \right\}$$

Pf:

$$|x_n - \xi| = |g(x_{n-1}) - g(\xi)| \stackrel{\text{MVT}}{=} |g'(\xi)| \cdot |x_{n-1} - \xi| \leq m |x_{n-1} - \xi|$$

for some $\xi \in (\xi, x_{n-1})$

MVT

$$\Rightarrow |x_n - \xi| \leq m^n |x_0 - \xi|$$

$$|x_{n+k+1} - x_{n+k}| = |g(x_{n+k}) - g(x_{n+k-1})| \leq m |x_{n+k} - x_{n+k-1}| \leq m^k |x_{n+1} - x_n|$$

②)

$$\Rightarrow |x_{n+1} - \xi| \leq |x_{n+2} - x_{n+1}| + |x_{n+3} - x_{n+2}| + \dots \leq (m + m^2 + \dots) |x_{n+1} - x_n| = \frac{m}{1-m} |x_{n+1} - x_n| \leq \frac{m^{n+1}}{1-m} (x_1 - x_0)$$

Thm (Speed of convergence)

In the neighborhood of ξ assuming it is p -differentiable $\dot{}$

$$g'(\xi) = g''(\xi) = \dots = g^{(p-1)}(\xi) = 0$$
$$g^{(p)}(\xi) \neq 0$$

then order of convergence is p

Pf: Write Taylor series (since all derivatives to p are zero)

$$x_{n+1} = g(x_n) \approx \xi + \frac{1}{p!} g^{(p)}(\xi) (x_n - \xi)^p$$

↓
MVT

$$\frac{|x_{n+1} - \xi|}{|x_n - \xi|^p} = \frac{1}{p!} |g^{(p)}(\xi)|$$

fzero: MATLAB function

python: root_scalar, brentq

Julia: Roots.jl find_zero

Systems of nonlinear equations

$$f_1(x_1, \dots, x_n) = 0$$

$$f_2(x_1, \dots, x_n) = 0$$

⋮

$$f_n(x_1, \dots, x_n) = 0$$

$$\underline{f} := (f_1 \dots f_n) : \mathbb{R}^n \rightarrow \mathbb{R}^n$$

We write $\underline{f}(\underline{x}) = 0$

Newton (tangent method)

Jacobi (fixed point)

Newton iteration

$$\underline{x}^{(k+1)} = \underline{x}^{(k)} - \underline{J}_f(\underline{x}^{(k)})^{-1} \underline{f}(\underline{x}^{(k)})$$

$$\underline{J}_f(\underline{x}) = \begin{bmatrix} \frac{\partial f_1}{\partial x_1} & \dots & \frac{\partial f_1}{\partial x_n} \\ \vdots & & \vdots \\ \frac{\partial f_n}{\partial x_1} & \dots & \frac{\partial f_n}{\partial x_n} \end{bmatrix}(\underline{x})$$



Jacobi matrix

in practice we don't compute the inverse but solve the system

$$\begin{aligned} \text{linear system} \rightarrow J_f(\underline{x}^{(r)}) \Delta \underline{x}^{(r)} &= -f(\underline{x}^{(r)}) \\ \underline{x}^{(r+1)} &= \underline{x}^{(r)} + \Delta \underline{x}^{(r)} \end{aligned}$$

Derivation

① Write function via Taylor series

$$f_i(\underline{x} + \Delta \underline{x}) = f_i(\underline{x}) + \sum_{k=1}^n \frac{\partial f_i}{\partial x_k}(\underline{x}) \Delta x_k + \dots$$

② Drop higher order terms & set

$$f_i(\underline{x} + \Delta \underline{x}) = 0$$

③ We get the above system

Jacobi iteration

$$\begin{aligned} \text{① } f(\underline{x}) = 0 &\Leftrightarrow g(\underline{x}) = \underline{x} \\ &g: \mathbb{R}^n \rightarrow \mathbb{R}^n \end{aligned}$$

② Choose $\underline{x}^{(0)} \in \mathbb{R}^n$ (initial guess)

$$\text{③ } \underline{x}^{(r+1)} = g(\underline{x}^{(r)})$$

Thm (Jacobi iteration)

Given $g: \mathbb{R}^n \rightarrow \mathbb{R}^n$ on $\Omega \subset \mathbb{R}^n$

① $g(\Omega) \subset \Omega$ (closed on the set)

② $\|g(\underline{x}) - g(\underline{y})\| \leq m \|\underline{x} - \underline{y}\| \quad \forall \underline{x}, \underline{y} \in \Omega$ for some
(contracting map) $m \in [0, 1]$

Then $g(x) = x$ has a unique solution ξ

on Ω ; $\underline{x}^{(r+1)}$ converges to ξ for any $\underline{x}^{(0)}$

$$\text{Also } \|\underline{x}^{(r+1)} - \xi\| \leq \frac{m}{1-m} \|\underline{x}^{(1)} - \underline{x}^{(0)}\|$$

Thm (2nd convergence theorem for Jacobi iteration)

Let $g: \mathbb{R}^n \rightarrow \mathbb{R}^n$ be continuously differentiable at the fixed point ξ ; let $\|J_g(\xi)\| \leq 1$ then there exists a closed neighborhood

$\Omega \subseteq \mathbb{R}^n$ of ξ so that $\underline{x}^{(r+1)} \rightarrow \xi$

Note : the condition is different

Quasi-Newton methods

Broyden

- Newton: every step requires $O(n^2)$ partial derivatives & $O(n^3)$ operations for solving the linear system

- Broyden method: Like secant method avoids computing differentials

Let B_r approx. $J_f(x^{(r)})$

① Solve $B_r \Delta x^{(r)} = -f(x^{(r)})$

② $x^{(r+1)} = x^{(r)} + \Delta x^{(r)}$

③ Determine B_{r+1}

Broyden: Satisfies

$$B_{r+1}(x^{(r+1)} - x^{(r)}) = f(x^{(r+1)}) - f(x^{(r)})$$

∴ minimizes $\|B_{r+1} - B_r\|_2$

So we are looking for

$$\Delta B_r = B_{r+1} - B_r \quad w.l$$

$$\min \|\Delta B_r\|_2$$

$$\text{s.t. } \Delta B_r \Delta x^{(r)} = f(x^{(r+1)})$$

Claim: Jacobian matrix approx in Broyden method.

$$B_{r+1} = B_r + \frac{f(x^{(r+1)}) (\Delta x^{(r)})^T}{\|(\Delta x^{(r)})\|_2^2}$$

Pf: ① $B_r \Delta x^{(r)} = -f(x^{(r)})$

② $B_{r+1} \Delta x^{(r)} = f(x^{(r+1)}) - f(x^{(r)})$

combine ① & ②

o.l $(B_{r+1} - B_r) \Delta x^{(r)} = f(x^{(r+1)})$

so $B_{r+1} - B_r = \frac{f(x^{(r+1)}) (\Delta x^{(r)})^T}{\|\Delta x^{(r)}\|_2^2}$

\Rightarrow every matrix satisfying o.l \Rightarrow must have at least this norm. $\sim \frac{\|f(x^{(r+1)})\|_2}{\|\Delta x^{(r)}\|_2}$

Variational methods

If we have a good initial approximation

→ we can look for a local minimum

Claim

$$f(x) = 0$$

$$g: \mathbb{R}^n \rightarrow \mathbb{R} \quad g(x) = \|f(x)\|^2 = \sum_i f_i^2(x)$$

Q. if g is C^2 how can we find extremal points?

Finding local extremal points

Minimum of a function can be found iteratively in some direction v_r

$$x^{(r+1)} = x^{(r)} + \lambda_r v_r$$

where λ_r is some real number

$$g(x^{(r+1)}) < g(x^{(r)})$$

General descent: any direction not orthogonal to ∇g

Fastest descent: $-\nabla g = v_r$

Coordinate descent: $e_1, e_2, e_3 \dots e_n$

For solving step size

$$q(\lambda) = g(\underline{x}^{(k)} + \lambda v_r)$$

① gradient descent $q'(\lambda) = 0$

② tangent descent find intersection
tangent to $y = q(\lambda)$
at $\lambda = 0$ w/ x -axis

③ Quadratic parabolic descent

Use tangent to determine α

Fit a parabola to $(0, q(0)), (\frac{\alpha}{2}, q(\frac{\alpha}{2})),$

$(\alpha, q(\alpha))$ & take the minimum

Using methods for finding roots for finding extremal points

Claim: Local extremal points of $g(x)$ solve the system

$$\nabla g(x) = \left(\frac{\partial g(x)}{\partial x_1}, \dots, \frac{\partial g(x)}{\partial x_n} \right) = 0$$

Type of extremal point - depends on Hessian matrix

$$H_g(x) = \begin{bmatrix} \frac{\partial^2 g(x)}{\partial x_1^2} & \dots & \frac{\partial^2 g(x)}{\partial x_1 \partial x_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial^2 g(x)}{\partial x_n \partial x_1} & \dots & \frac{\partial^2 g(x)}{\partial x_n^2} \end{bmatrix}$$

if $H_g(x)$ is positive definite \Rightarrow it is a local minimum
negative definite \Rightarrow maximum

Gauss - Newton iteration

Solving overdetermined system

$$f: \mathbb{R}^n \rightarrow \mathbb{R}^m \quad f(x) = (0, \dots, 0)$$

m non linear equations in n unknowns
 $m > n$

Since generally no solution exists, find solution s.t.

$$\|f(x)\|^2 = \min \{ \|f(x)\|^2 \}$$

Gauss-Newton

$$\underline{x}^{(r+1)} = \underline{x}^{(r)} - \underbrace{J_r(\underline{x}^{(r)})^+}_{\text{pseudo inverse}} f(\underline{x}^{(r)})$$

rather than inverse

At each step solve

$\underline{x}^{(r+1)}$ is the "solution" of

$$J_r(\underline{x}^{(r)}) \Delta \underline{x}^{(r)} = f(\underline{x}^{(r)})$$

no convergence \Rightarrow even when there is, it is only local convergence

Example:

$$(x_i, y_i) \in \mathbb{R}^2 \quad i=1, \dots, m$$

Find a function $f(x, a, b) = ae^{bx}$

$$F: \mathbb{R}^2 \rightarrow \mathbb{R}^m$$

$$F(a, b) = (y_1 - ae^{bx_1}, \dots, y_m - ae^{bx_m})$$

$$DF(a, b) = \begin{bmatrix} -e^{bx_1} & ax_1 e^{bx_1} \\ \vdots & \vdots \\ -e^{bx_m} & ax_m e^{bx_m} \end{bmatrix}$$

$$\begin{bmatrix} a_{r+1} \\ b_{r+1} \end{bmatrix} = \begin{bmatrix} a_r \\ b_r \end{bmatrix} - DF^+(a_r, b_r) F(a_r, b_r)^T$$

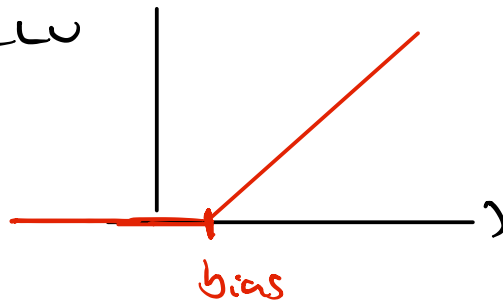
What about non-differentiable functions?

① Automatic differentiation - backprop

② Almost differentiable

a) - max (maxpooling) $\max(x)$

b) - ReLU



②

b) Just use the fact that we never hit the corner point (always evaluate $f'(x_k)$)

a) Say multiple entries are equal

$$y = (x_1, \dots, x_n)$$

$$\nabla y = (0 \dots 0 \mid \dots \mid 0) \text{ if unique}$$

otherwise

$$\nabla y = (0 \dots 0 \mid \frac{1}{k} \dots \frac{1}{k} \mid 0 \dots 0) \quad \frac{1}{k} \text{ for } k \text{ max values}$$

or choose randomly

\Rightarrow practical solution

① Idea (will cover later)

\Rightarrow write derivatives of basic function

\Rightarrow use chain rule to construct derivative

$$\frac{\partial}{\partial x} f(g(x)) = f'(g(x)) g'(x)$$

forward $x \mapsto (x, 1)$ $\sin(x) \mapsto (\sin(x), \cos(x))$
 $x^2 \mapsto (x^2, 2x)$

backward

$$\frac{\partial L}{\partial x} = \frac{\partial L}{\partial y} \cdot \frac{\partial y}{\partial x}$$

Ex $f(x) = (x^2 + 3x) \sin x$ $a = x^2$
 $x \mapsto (x, 1)$ $b = 3x$
 $a: x^2 \mapsto (x^2, 2x)$ $c = a + b$
 $b: b \mapsto (3x, 3)$ $d = \sin x$
 $d: \sin x \mapsto (\sin x, \cos x)$ $e = c \cdot d$

$c \mapsto (a + b, a' + b')$
 $(x^2 + 3x, 2x + 3)$

$e: c \cdot d \mapsto (f(x), c' \cdot d + d' \cdot c)$

Dual numbers \Rightarrow usually at a point

$$x = a + \epsilon \Rightarrow x = a + f'(x) \epsilon$$

Start at $(y, \epsilon) \rightarrow f(y) = (f(y), f'(y) \epsilon)$

Multiply 2 dual numbers

$$(a+b\varepsilon)(c+d\varepsilon) = \underline{ac} + \underline{(ad+bc)}\varepsilon + bd\varepsilon^2$$

↓
set to
0