

Digital forensics

Andrej Brodnik

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Operating system Unix

Chapter 18

- Development through history: System V, HP-UX, BSD, ...
- Open source versions that appeared later:
 - Linux: RedHat, SUSE, Ubuntu, ...
 - BSD: FreeBSD, OpenBSD, NetBSD

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File system Hierarchy Standard

- File system Hierarchy Standard – FHS
(<http://www.pathname.com/fhs/pub/fhs-2.3.html>)
- Linux Foundation took over the work
(<http://www.linuxfoundation.org/collaborate/workgroups/lfs/fhs>)
- Mostly formalization of the BSD file system

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Root directory

- */boot* : Static files of the boot loader
- */dev* : Device files
- */etc* : Host-specific system configuration
 - */etc/opt* : Configuration files for */opt*
 - */etc/X11* : Configuration for the X Window System (optional)
 - */etc/sgml* : Configuration files for SGML (optional)
 - */etc/xml* : Configuration files for XML (optional)
- */bin* : Essential user command binaries (for use by all users)
- */sbin* : System binaries
- */lib* : Essential shared libraries and kernel modules
- */lib<qual>* : Alternate format essential shared libraries (optional)

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Root directory

- */home* : User home directories (optional)
- */root* : Home directory for the root user (optional)
- */media* : Mount point for removable media
- */mnt* : Mount point for a temporarily mounted filesystem
- */opt* : Add-on application software packages
- */srv* : Data for services provided by this system
- */tmp* : Temporary files
- */usr*, */var* : Separate hierarchies

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/usr directory

- Contains read-only files
- Used simultaneously by different systems
- Doesn't contain files that are specific to a particular system
- Exception: */usr/local*, which is the local directory of a particular system

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/var directory

- Contains files that change over time
 - Postal and print queues
 - Logging
 - Data (databases etc)
 - Temporary files

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System files

- Operating system is designed so that system files are user-friendly → regular text files
 - Configuration files: hosts, syslog.conf, ...
 - Usually in the directory etc (/etc, /usr/local/etc, /opt/etc, ...)
 - Logging : mail, cups, ...
 - Usually in the directory log (/var/log, /usr/local/var/log, /opt/var/log)

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Configuration files

```
# $FreeBSD: release/9.0.0/etc/snmpd.config.216595 2010-12-20 17:28:15Z syrinx $
#
# Example configuration file for bsnmppd(1).
#
#
# Set some common variables
#
location := "Room 200"
contact := "sysmaster@example.com"
system := 1 # FreeBSD
traphost := localhost
trapport := 162

#
# Set the SNMP engine ID.
# The snmpEngineID object required from the SNMPv3 Framework. If not explicitly set via
# this configuration file, an ID is assigned based on the value of the
# kern.hostid variable
# engine := 0x80:0x10:0x08:0x10:0x80:0x25
# snmpEngineID = $(engine)
```

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Logging

Mar 8 00:00:00 svarun newsyslog[85254]: logfile turned over
Mar 8 00:00:12 svarun postfix/smtpd[85247]: connect from S0106c0c1c0ddffcf.vf.shawcable.net[70.69.32.154]
Mar 8 00:00:12 svarun postfix/smtpd[85247]: NOQUEUE: reject: RCPT from S0106c0c1c0ddffcf.vf.shawcable.net[70.69.32.154]: 554 5.7.1 Service unavailable; Client host [70.69.32.154] blocked using bl.spamcop.net; Blocked - see <http://www.spamcop.net/bl.shtml?70.69.32.154>; from=sunscrupulousnessw2@deltamar.net> to=<xxxx@brodnik.org> proto=ESMTP helo=<deltamar.net>
Mar 8 00:00:12 svarun postfix/smtpd[85247]: lost connection after DATA from S0106c0c1c0ddffcf.vf.shawcable.net[70.69.32.154]

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Data storage and hiding

- Simplified organization of the disk with FAT file system

The diagram illustrates the layout of a disk with two partitions, Partition #1 and Partition #2. Each partition starts with a boot sector, followed by a primary FAT table, a secondary FAT table, and a root directory. The main area of each partition is the file space and unallocated space, which is divided into volume slack and unused disk space. The entire disk is labeled as Sector 1 (MBR).

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File systems

- We have directories and index nodes(*inode*)
- Inode has a similar function to FAT and MFT at the same time
- Directory is just a special file type
 - We have more special files: links, pipe, socket ...

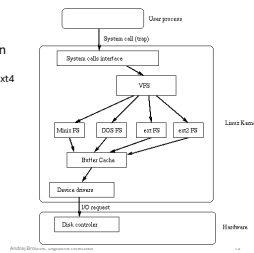
The diagram shows a root directory containing entries like '163841', 'var', '212963', and '229377'. An arrow points from the '229377' entry to an inode structure. The inode structure contains fields for owner/groupID, permission, file type, time stamps, reference count, file size in bytes, and data blocks #s. An arrow points from the 'data blocks #s' field to a set of data blocks.

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File systems

- The oldest: Unix File System – UFS
- More recent and used in Linux: ext2 in ext3
 - There are also ext and ext4
- There are a number of other file systems



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Time in the Unix operating system

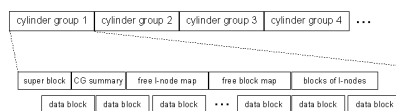
- Time is measured in seconds
 - If time is stored as a 32-bit number, there will be an overflow on Tuesday, January 19th 2038 at 03:14:07 UTC – Y2K38 problem
- UTC – *Coordinate Universal Time*: a harmonized definition of time that takes into account leap years, leap seconds, ...
 - The last leap second occurred on 31st January 2016
 - harmonized time between several atomic hours
 - one of the successors of GMT

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UFS file systems

- Defined when VFS was introduced in BSD4.2
- Used in *BSD systems
- Later used in Solaris OS

vir: *Solaris Internals, The UFS File System*, Updated by Frank Batschulat, Shawn Debnath, Sarah Jelinek, Dworkin Muller, and Karen Rochford



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UFS – index node

```

struct dinode {
    u_int16_t    d1_mode;      /* 0: IFMT, permissions; see below. */
    int16_t      d1_link;      /* 2: File link count. */
    u_int16_t    d1_dauid[2];  /* 4: File's old user and group ids. */
    int32_t      d1_inumber;   /* 4: Lfs: inode number. */
    int32_t      d1_size;      /* 8: File byte count. */
    int32_t      d1_atime;     /* 16: Last access time. */
    int32_t      d1_mtime;     /* 20: Last access time. */
    int32_t      d1_ctime;     /* 24: Last modified time. */
    int32_t      d1_chtime;    /* 32: Last inode change time. */
    ufs_daddr_t  d1_db[NDAADDR]; /* 40: Direct disk blocks. */
    ufs_daddr_t  d1_ib[NDAADDR]; /* 88: Indirect disk blocks. */
    u_int32_t     d1_flags;     /* 100: Status flags (chflags). */
    int32_t      d1_bcount;    /* 108: Blocks currently held. */
    u_int32_t     d1_btag;     /* 112: File owner. */
    int32_t      d1_spare[2];  /* 120: Reserved; currently unused */
};

```

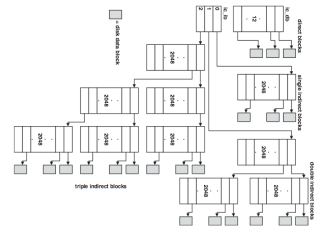
ufs/dinode.h

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UFS – file systems



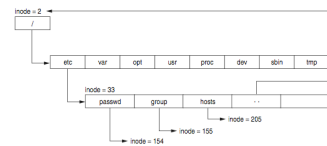
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UFS – directory file

- a special file that consists of parts of the directory
- System V had a predefined file size
- The root directory is described in inode 2
- Each directory has a special entry that notes where its parent is



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UFS – directory entry

```
#define MAXNAMLEN 255
struct direct {
    u_int32_t d_ino; /* inode number of entry */
    u_int16_t d_reclen; /* length of this record */
    u_int8_t d_type; /* file type, see below */
    u_int8_t d_namlen; /* length of string in d_name */
    char d_name[MAXNAMLEN + 1];
};
```

ufs/dir.h

- Challenge: what is the record reclen intended for? Can this be used to hide data?
- Challenge: what is ACL? How is it implemented in UFS?

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UFS – superblock

- Superblock stores the description of the cylinder group's configuration
- Scattered around the disc - at the beginning of each group of cylinders
- To save the configuration – if one record is lost
- dumpfs** tool
- Challenge: find the structure of superblock. How do we know that we are dealing with the UFS file system? Where is it written? Read the superblock from your unix file system and find out for which file system it is.

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File system ext2

- The basic structure is similar to that of UFS
- Instead of groups of cylinders, we are talking about block groups
- Directories and index nodes - like in UFS

The diagram illustrates the layout of an ext2 file system. It starts with Sector 1 (MBR) containing the Partition table. This leads to Partition #1, which contains the Superblock1, Group descriptor table, Block bitmap, Inode bitmap, File system and unallocated space, Superblock1 (backup), Group descriptor table, Block bitmap, Inode bitmap, File system and unallocated space, and Superblock2. Partition #2 follows. Below the diagram, Group #1 and Group #2 are indicated, showing the distribution of data across different block groups.

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File system ext2

- Tool for viewing disks: Linux Disk Editor (LDE)
(<http://lde.sourceforge.net/>)

```
lde v2.6.0 : ext2 : /dev/hdd2
Inode:      2 (0x00000002)  Block:      0 (0x00000000)

0x00000002: dswr-xr-x 21 4096 .
0x00000002: dswr-xr-x 21 4096 ..
0x00000008: dswr-xr-x 2 16384 lost+found
0x00008001: dswr-xr-x 2 4096 boot
0x00010001: dswr-xr-x 17 77824 dev
0x00020001: dswr-xr-x 2 4096 proc
0x0000000C: rw-r--r-- 1 0 .autofsck
0x00028001: dswr-xr-x 17 4096 var
0x00034001: dswr-rw-rw- 8 4096 tmp
0x00038001: dswr-xr-x 40 4096 etc
0x00048001: dswr-xr-x 15 4096 usr
0x00058003: dswr-xr-x 2 4096 bin
0x00064003: dswr-xr-x 3 4096 home
0x00064003: dswr-xr-x 2 4096 initrd
0x00065003: dswr-xr-x 7 4096 lib
0x00066003: dswr-xr-x 4 4096 mnt
0x00066C03: dswr-xr-x 2 4096 opt
0x00067003: dswr-xr-x 7 4096 root
0x00070003: dswr-xr-x 2 4096/sbin
0x00044C04C: dswr-xr-x 2 4096 misc
0x00000021: dswr-xr-x 4 4096 ai
```

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File system ext2

```
lde v2.6.0 : ext2 : /dev/hdd2
Inode:      229505 (0x0038081)  Block:      0 (0x00000000)

-rw-r--r-- 1 root root 1186 Tue Sep 24 08:57:40 2002

TYPE: regular file LINKS: 1 DIRECT BLOCKS=0x000703F9
MODE: \0644 FLAGS: \10
UID: 00000(root) GID: 00000(root)
SIZE: 1186 STIM(BLOCKS): 8

ACCESS TIME: Tue Nov 26 11:10:18 2002
CREATION TIME: Tue Sep 24 08:57:40 2002
MODIFICATION TIME: Tue Sep 24 08:57:40 2002
DELETION TIME: Wed Dec 31 19:00:00 1969

INDIRECT BLOCK=
2x INDIRECT BLOCK=
3x INDIRECT BLOCK=
```

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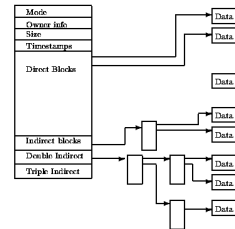
ext2 – index nodes

```
struct ext2_inode {
    __u16 i_mode; /* 0: File mode */
    __u16 i_uid; /* 2: Owner Uid */
    __u32 i_size; /* 4: Size in bytes */
    __u32 i_atime; /* 8: Access time */
    __u32 i_ctime; /* 12: Creation time */
    __u32 i_mtime; /* 16: Modification time */
    __u32 i_dtime; /* 20: Deletion time */
    __u16 i_gid; /* 24: Group id */
    __u16 i_links_count; /* 28: Links count */
    __u32 i_block; /* 32: Blocks count */
    __u32 i_flags; /* 36: File flags */
    __u32 i_i_reserved1; /* 36: OS dependent 1 */
    __u32 i_block[XEXT2_N_BLOCKS]; /* 40: Pointers to blocks */
    __u32 i_generation; /* 104: File version (for NFS) */
    __u32 i_file_acl; /* 104: File ACL */
    __u32 i_dir_acl; /* 108: Directory ACL */
    __u32 i_addr; /* 112: Fragment address */
    __u8 i_frag; /* 116: Fragment number */
    __u8 i_fsize; /* 117: Fragment size */
    __u32 i_i_reserved2[2]; /* 120: OS dependent 2 */
};
```

ext2fs/ext2_fs.h

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ext2 – index nodes



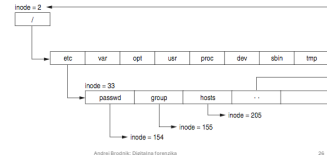
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Directory file

- A special file that consists of parts of the directory
- System V had a predefined file size
- The root directory is described in inode 2
- Each directory has a special entry .. that tells where the parent is



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ext2 – Directory entry

```
#define EXT2FS_MAXNAMLEN 255
struct ext2fs_direct {
    u_int32_t e2d_ino; /* inode number of entry */
    u_int16_t e2d_reclen; /* length of this record */
    u_int8_t e2d_namlen; /* length of string in d_name */
    u_int8_t e2d_type; /* file type */
    char e2d_name[EXT2FS_MAXNAMLEN]; /* name with length <=
EXT2FS_MAXNAMLEN */
};
ext2fs/ext2fs_dir.h
```

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ext2 – superblock

- The superblock stores the description of the block group configuration
- Scattered across the disk - at the beginning of each block of blocks
 - to save the configuration if one record is lost
- Tool **dumpfs**

• Izziv: poiščite strukturo nadbloka ext2. Primerjajte jo s strukturo UFS superbloka.

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File system ext3

- Author Stephen Tweedie 1999 / 2000 / 2001
- The basic structure is the same as for the ext2 file system
 - Split into blocks of blocks including a superblock
 - Directories and index nodes
 - Keeping track of the disk
- The option of saving the log structure is added
- The basic OS Linux file system

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Journals ext3

- Journals contain records of all changes to the file system
- Journal's structure allows for three types of journals:
 - comprehensive journal: saves everything; both metadata and content - the most secure
 - ordered: only metadata is stored but only after a successful operation - medium safe
 - writeback: similar to sequential, saving log records at the same time as actual records - least secure

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Journals ext3

- Journal is a sequential file
- Records are stored in front of the first group of blocks
- The log group is similar to the block group:
 - Journal superblock
 - Transaction descriptions

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Journal ext3

- The transaction description contains three types of blocks:
 - Descriptor block: start of a transaction
 - Metadata blocks: transaction descriptions
 - Commit block: completion of the transaction
 - Revoke block: if an error occurs and contains a list of blocks in the file system that need to be reinstalled (restored)
- All (including superblock) start with a magical number:
 - JFS_DESCRIPTOR_BLOCK1
 - JFS_COMMIT_BLOCK2
 - JFS_SUPERBLOCK_V13
 - JFS_SUPERBLOCK_V24
 - JFS_REVOKE_BLOCK5

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Journal ext3

- Challenge: Consider the structure of a superblock (e.g. <http://linuxsoftware.co.nz/wiki/ext3>) . Get a block from your file system and comment on its contents.
- Challenge: How do we restored a deleted file in ext2 and in ext3? What about in UFS?

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File systems

- There are other file systems:
 - reiserFS, XFS, gfs, afs, ext4, HSM, ...
- *Challenge:* Make a similar analysis for the mentioned systems as we did for UFS and ext.
- *Challenge:* Compare the described file systems – in which can we hide data?
- *Challenge:* Prepare a file system for your colleague and he must figure out which one it is.

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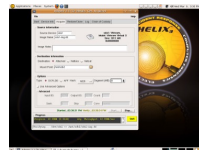
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Forensic sources

- We use stand-alone operating systems to analyse the disc image
- Example: Helix (Ubuntu)

- *Challenge:* Prepare the Helix CD and check what tools are already on it.
- *Challenge:* Find some other similar systems.



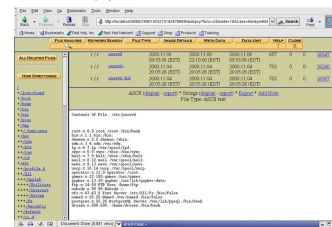
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Forensic sources

- Tool SleuthKit with Autopsy Forensic Browser



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Forensic sources

• Some interesting and rich references:

• B. Carter, *File system forensic analysis*. Addison-Wesley, 2005.

• Gregorio Narváez, *Taking advantage of Ext3 journaling file system in a forensic investigation*. SANS Institute, 2007.

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